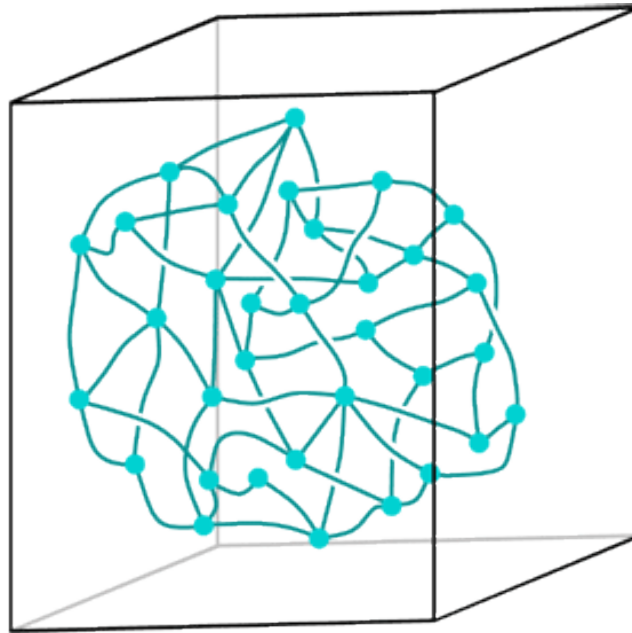
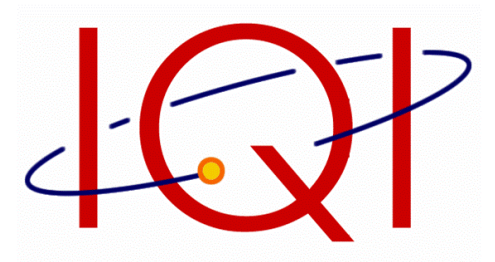


Permutational Quantum Computation

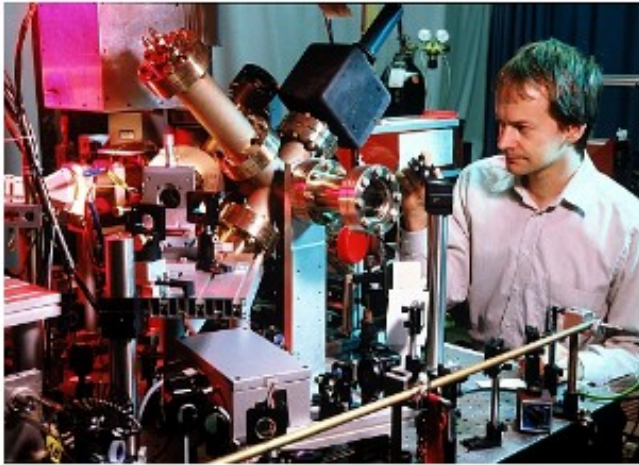


Stephen Jordan



Why formulate new models?

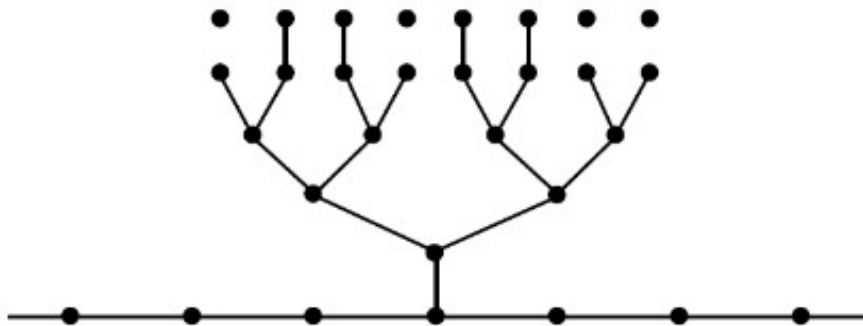
Implementation



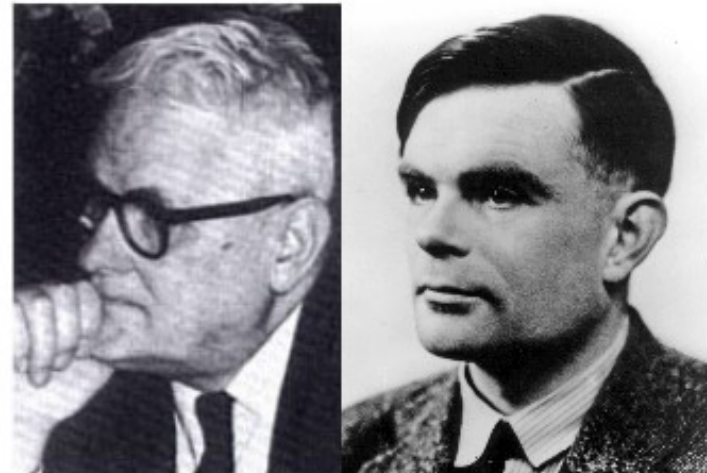
Complexity Theory



Quantum Algorithms



Church-Turing Thesis



- **what is the model?**
 - start with spins of known total angular momentum
 - permute the particles around
 - measure total angular momentum
 - direct analogue to topological quantum computation
- **what can it do?**
 - approximate irreps of S_n
 - approximate Ponzano-Regge invariant
 - give us a new complexity class?

Angular momentum of n spins

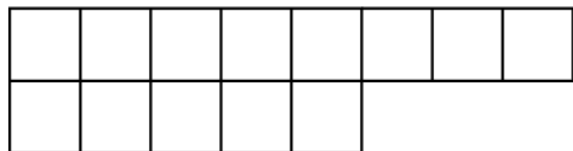
$$\vec{S}_j = \frac{1}{2} \begin{pmatrix} \sigma_x^{(j)} \\ \sigma_y^{(j)} \\ \sigma_z^{(j)} \end{pmatrix}$$

$$\vec{S} = \sum_{i=1}^n \vec{S}_i$$

$$S^2 = \vec{S} \cdot \vec{S}$$

$$S^2 |j\rangle = j(j+1) |j\rangle$$

- S^2 commutes with any permutation
- the eigenspaces of S^2 transform as irreducible representations of S_n
- The Young diagrams have two rows:



- The overhang is $2j$

- what about a basis for the representations?
- Example: 3 particles

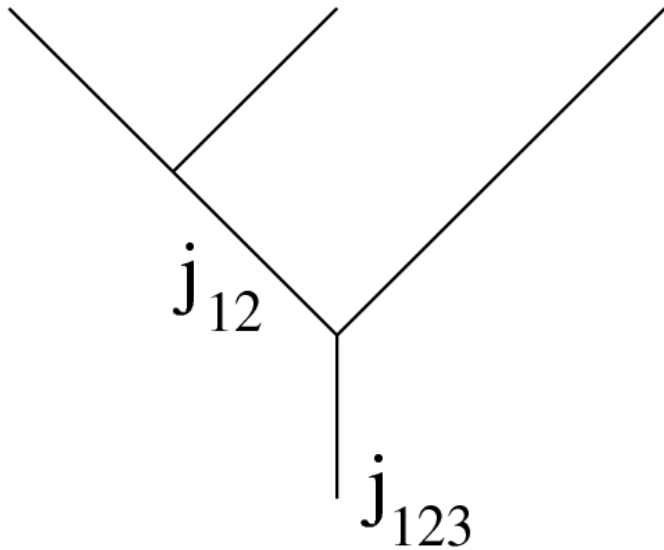
$$\left. \begin{array}{l} (\vec{S}_1 + \vec{S}_2 + \vec{S}_3)^2 \\ (\vec{S}_1 + \vec{S}_2)^2 \\ Z_1 + Z_2 + Z_3 \end{array} \right\} \text{complete set of} \\ \text{commuting} \\ \text{observables}$$

- How do the representations of S_3 look in this basis?
- $(\vec{S}_1 + \vec{S}_2 + \vec{S}_3)^2$ tells us which irrep
- $(\vec{S}_1 + \vec{S}_2)^2$ labels the basis states within an irrep
- $(Z_1 + Z_2 + Z_3)$ is an irrelevant degree of freedom

- We have a choice between bases:

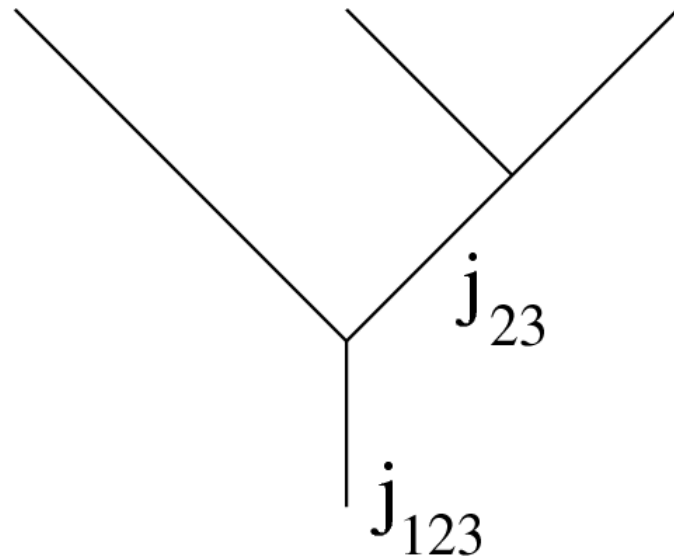
$$(\vec{S}_1 + \vec{S}_2 + \vec{S}_3)^2$$

$$(\vec{S}_1 + \vec{S}_2)^2$$



$$(\vec{S}_1 + \vec{S}_2 + \vec{S}_3)^2$$

$$(\vec{S}_2 + \vec{S}_3)^2$$



- For a given tree, different labellings correspond to orthogonal states

$$\left| \begin{array}{c} \diagup \quad \diagdown \\ \diagdown \quad \diagup \\ \text{a} \quad \text{b} \\ | \\ \text{a} \end{array} \right| \left| \begin{array}{c} \diagup \quad \diagdown \\ \diagdown \quad \diagup \\ \text{d} \quad \text{c} \\ | \\ \text{c} \end{array} \right| = \delta_{ac} \delta_{bd}$$

- Different trees are related by recoupling tensors

$$\left| \begin{array}{c} \text{a} \quad \text{b} \quad \text{c} \\ \diagdown \quad \diagup \\ \diagdown \quad \diagup \\ \text{d} \\ | \\ \text{e} \end{array} \right| = \sum_f \left[\begin{array}{ccc} \text{a} & \text{b} & \text{f} \\ \text{c} & \text{e} & \text{d} \end{array} \right] \left| \begin{array}{c} \text{a} \quad \text{b} \quad \text{c} \\ \diagdown \quad \diagup \\ \diagdown \quad \diagup \\ \text{f} \\ | \\ \text{e} \end{array} \right|$$

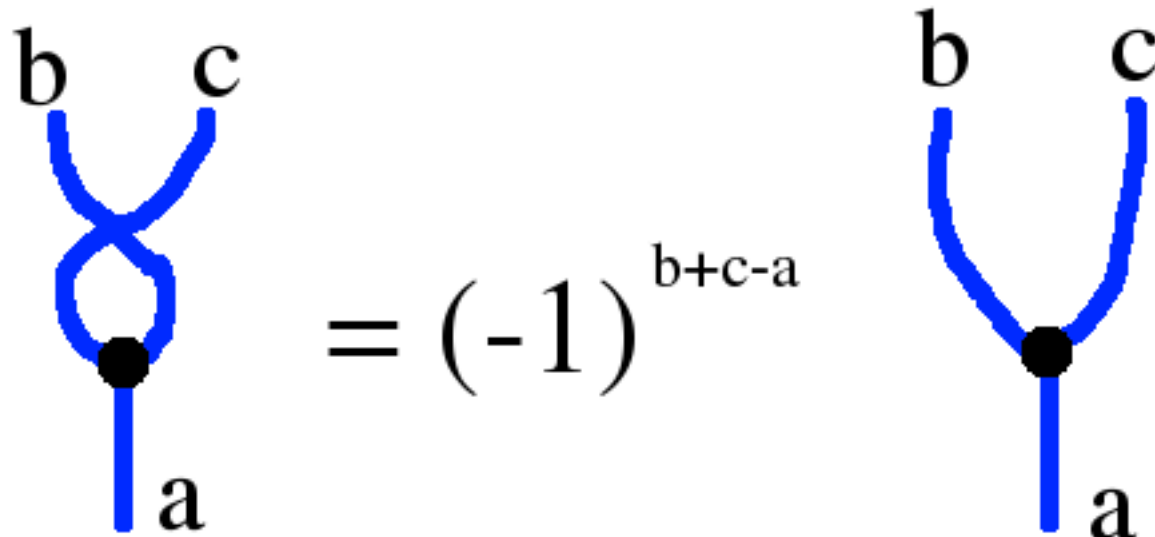
- The recoupling tensors are:

$$\begin{bmatrix} a & b & f \\ c & e & d \end{bmatrix} = (-1)^{a+b+c+f} \sqrt{(2d+1)(2f+1)} \left\{ \begin{matrix} a & b & f \\ c & e & d \end{matrix} \right\}$$

- The $6j$ symbols $\left\{ \begin{matrix} a & b & f \\ c & e & d \end{matrix} \right\}$ can be computed

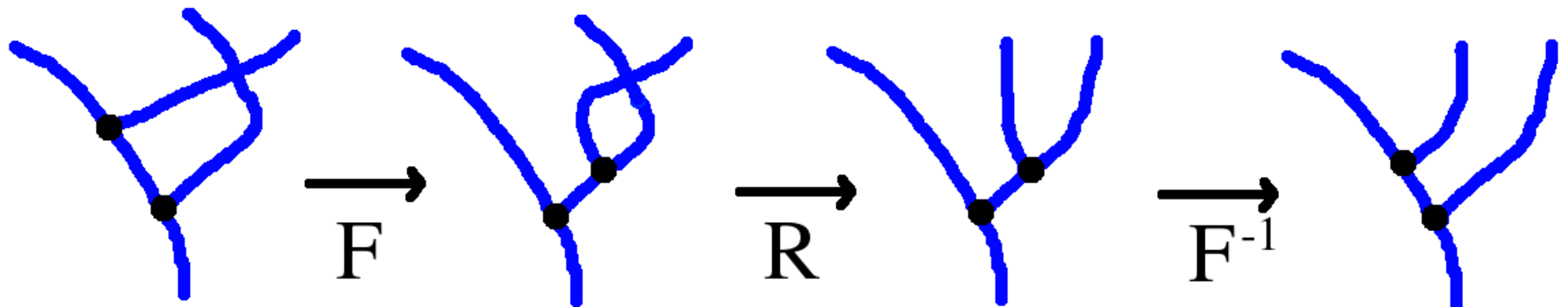
in $\text{poly}(a + b + c + d + e + f)$ time using the Racah formula.

- These states have this exchange symmetry:



$$\text{Diagram 1} = (-1)^{b+c-a} \text{Diagram 2}$$

- This plus recoupling tells us everything about permutation.



- We now have a model of computation:
 - 1) Prepare a basis state from some complete set of commuting angular momentum operators.
 - 2) Permute the qubits.
 - 3) Measure some other complete set of commuting angular momentum operators
- Variant: include phases

[See Also Marzuoli et al.]

Topological

Anyons

Braid (B_n)

Fuse

Braided Tensor
Category

Permutational

Spin-1/2

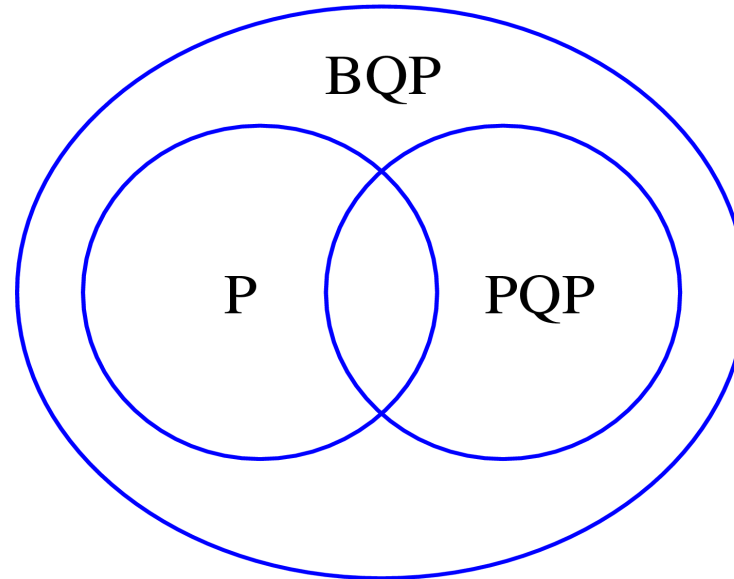
Permute (S_n)

Measure Angular
Momentum

Racah-Wigner
Tensor Category

How Powerful Is It?

- What I think:



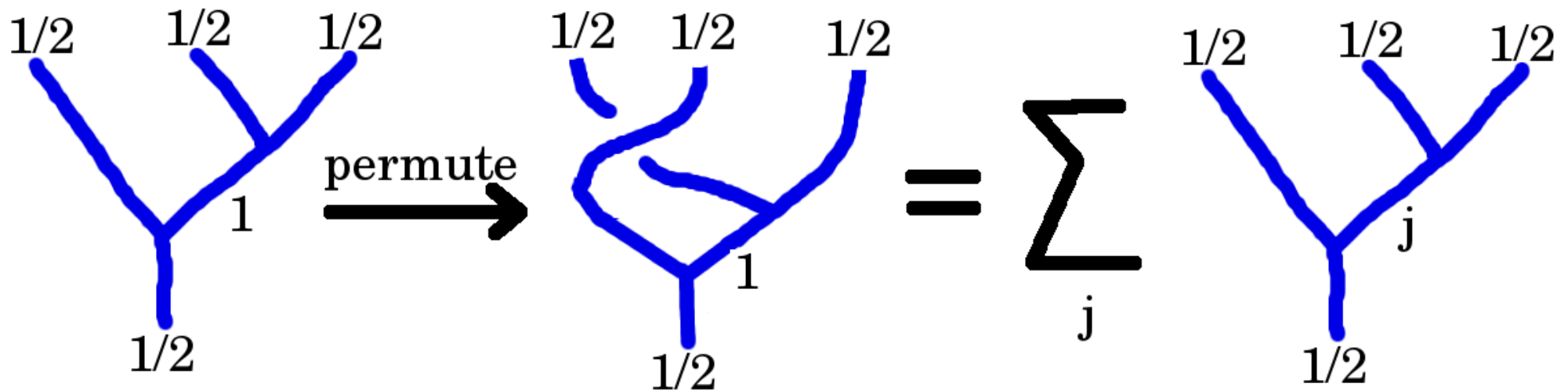
- What I know:

$$\text{PQP} \subset \text{BQP}$$

Can approximate irreps of S_n and simulate certain special cases of Ponzano-Regge.

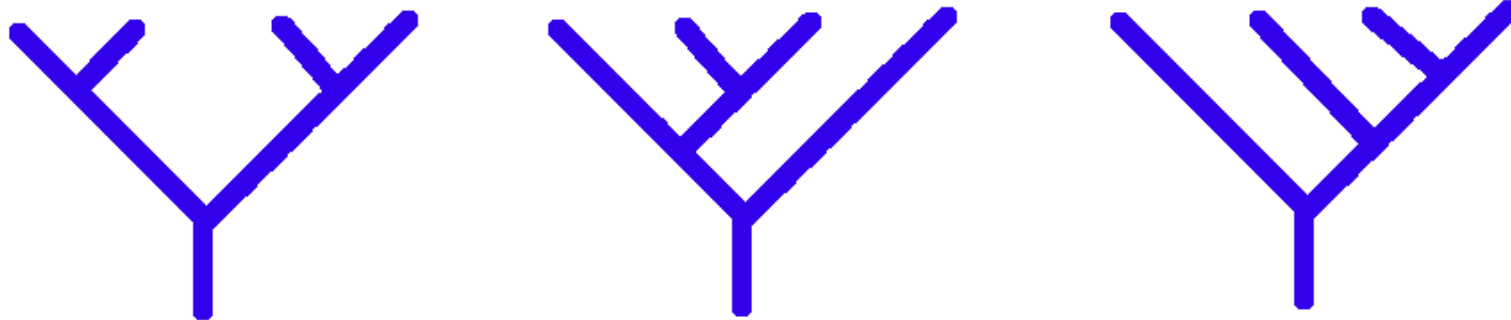
Algorithm for Symmetric Group

- permutation induces a linear transformation:

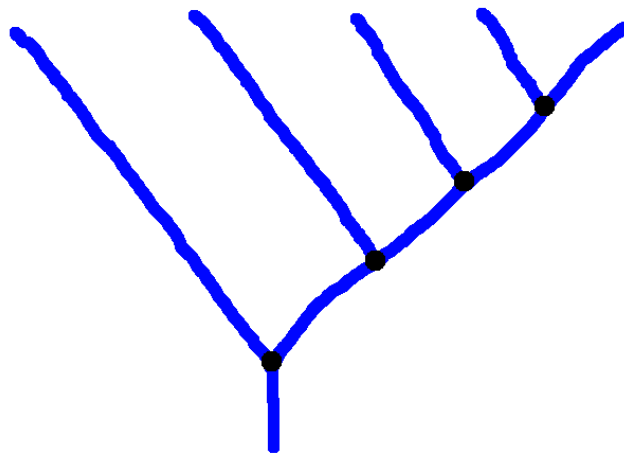


- This map from permutations to linear transformations is a representation of S_n

A choice of tree is a choice of basis.



If we choose this type of tree



then the representation of S_n is in Young's orthogonal form.

Symmetric-group-based methods in quantum chemistry

Jacek Karwowski

J. Phys. A: Math. Gen. 25 (1992) 3737-3747. Printed in the UK

An efficient algorithm for evaluating the standard Young-Yamanouchi orthogonal representation with two-column Young tableaux for symmetric groups

Wei Wu and Qianer Zhang

MATHEMATICS OF COMPUTATION
VOLUME 53, NUMBER 192
OCTOBER 1990, PAGES 705-722

COMPUTING IRREDUCIBLE REPRESENTATIONS OF FINITE GROUPS

LÁSZLÓ BABAI AND LAJOS RÓNYAI

ABSTRACT. We consider the bit-complexity of the problem stated in the title. Exact computations in algebraic number fields are performed symbolically. We present a polynomial-time algorithm to find a complete set of nonequivalent irreducible representations over the field of complex numbers of a finite group given by its multiplication table. In particular, it follows that some representative of each equivalence class of irreducible representations admits a

PROCEEDINGS OF THE
AMERICAN MATHEMATICAL SOCIETY
Volume 83, Number 2, October, 1981

A SIMPLIFICATION OF THE COMPUTATION OF THE NATURAL REPRESENTATION OF THE SYMMETRIC GROUP S_n

JOSEPH M. CLIFTON

INTERNATIONAL JOURNAL OF QUANTUM CHEMISTRY, VOL. 50, 55-67(1994)

The Orthogonal and the Natural Representation for Symmetric Groups

WEI WU* AND QIANER ZHANG

CHEMICAL PHYSICS LETTERS

1 April 1977

A RECURSIVE FORMULA FOR YOUNG'S ORTHOGONAL REPRESENTATION

Sten PETTRUP

Journal of Chemical Physics, The Technical University of Denmark, Lyngby, Denmark.

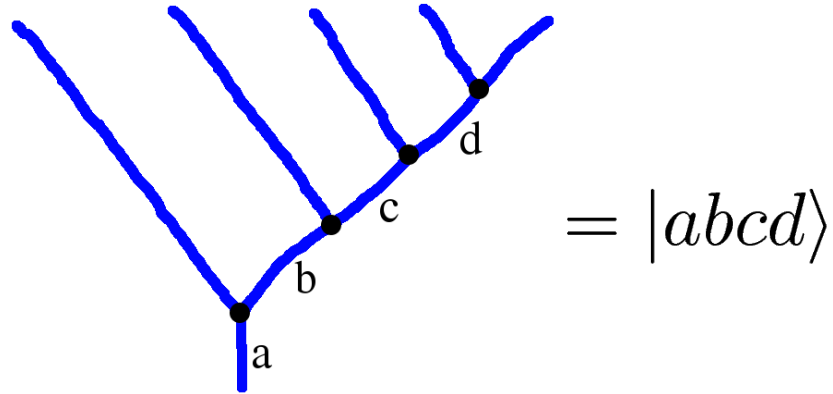
Matrices of Young's orthogonal representation can be generated by a recursive formula. The method is appropriate to computer application.

Systems can be treated by application of the theory of matrices as demonstrated by Kotani et al. [1]. The method is identical with the Yamanouchi method. The matrices can either be obtained from the recursive formula here $j = 1, 2, \dots, n$.

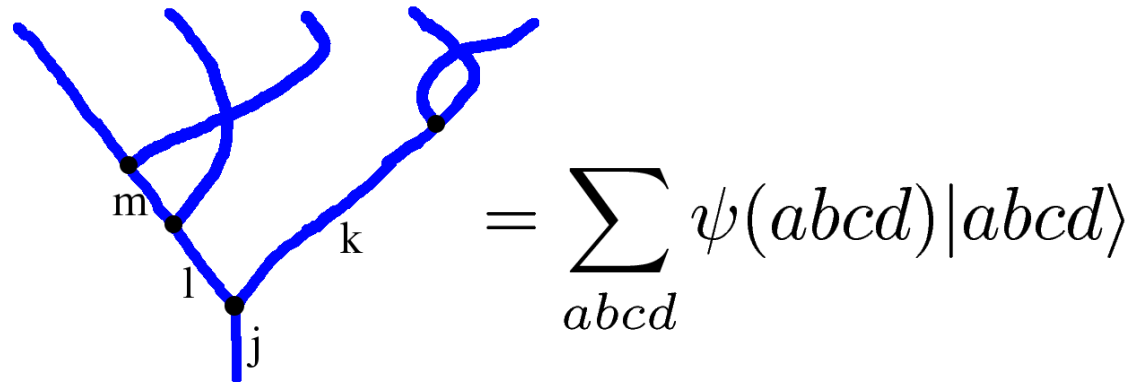
PQP \subset BQP

Proof Sketch:

work in this basis:

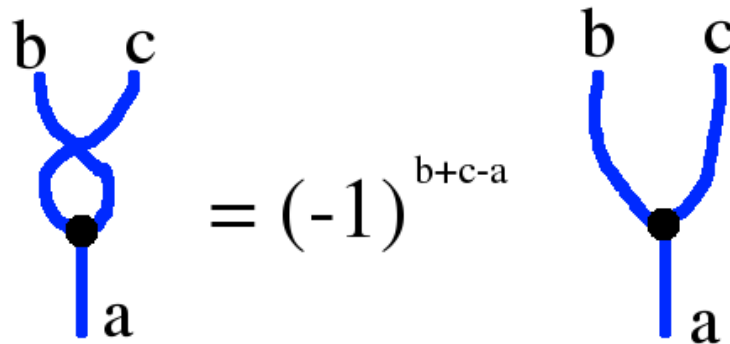


make any PQP state
by polynomially many
F and R moves



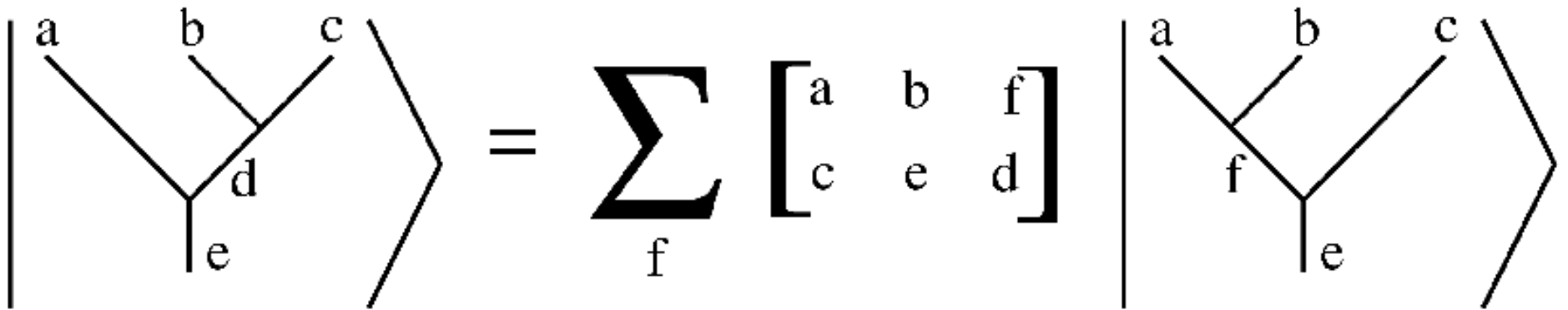
use Hadamard test

- R is easy to implement: just a phase



$$\begin{array}{c} b \\ \curvearrowright \\ \bullet \\ \curvearrowleft \\ c \\ | \\ a \end{array} = (-1)^{b+c-a} \begin{array}{c} b \\ | \\ \bullet \\ | \\ c \\ | \\ a \end{array}$$

- How about F?



$$\left| \begin{array}{c} a \\ \diagdown \\ \diagup \\ \diagdown \\ \diagup \\ e \\ | \\ e \end{array} \right. = \sum_f \left[\begin{array}{ccc} a & b & f \\ c & e & d \end{array} \right] \left| \begin{array}{c} a \\ \diagdown \\ \diagup \\ \diagdown \\ \diagup \\ e \\ | \\ e \end{array} \right.$$

- it is sparse
- we can efficiently compute the nonzero entries using the Racah formula

- We know how to implement any sparse row-computable **Hamiltonian**.
- From this we can implement any sparse row- and column-computable **unitary**.

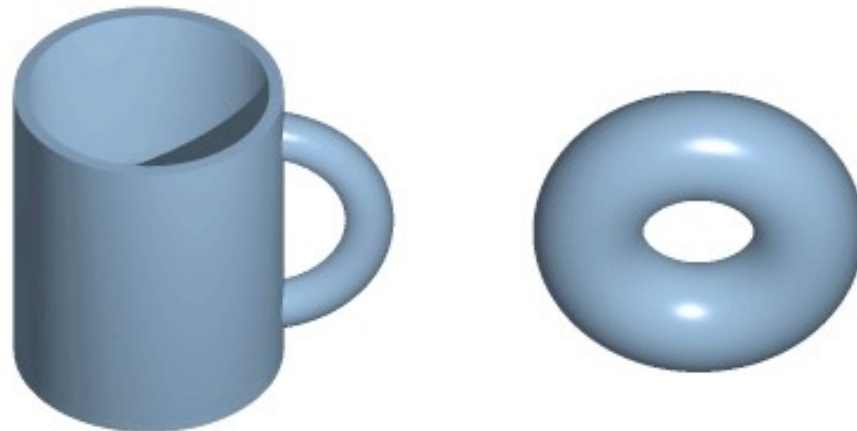
$$H = \begin{bmatrix} 0 & U \\ U^\dagger & 0 \end{bmatrix}$$

$$e^{iH\pi/2} = i \begin{bmatrix} 0 & U \\ U^\dagger & 0 \end{bmatrix} \quad \blacksquare \quad \text{End of Proof Sketch.}$$

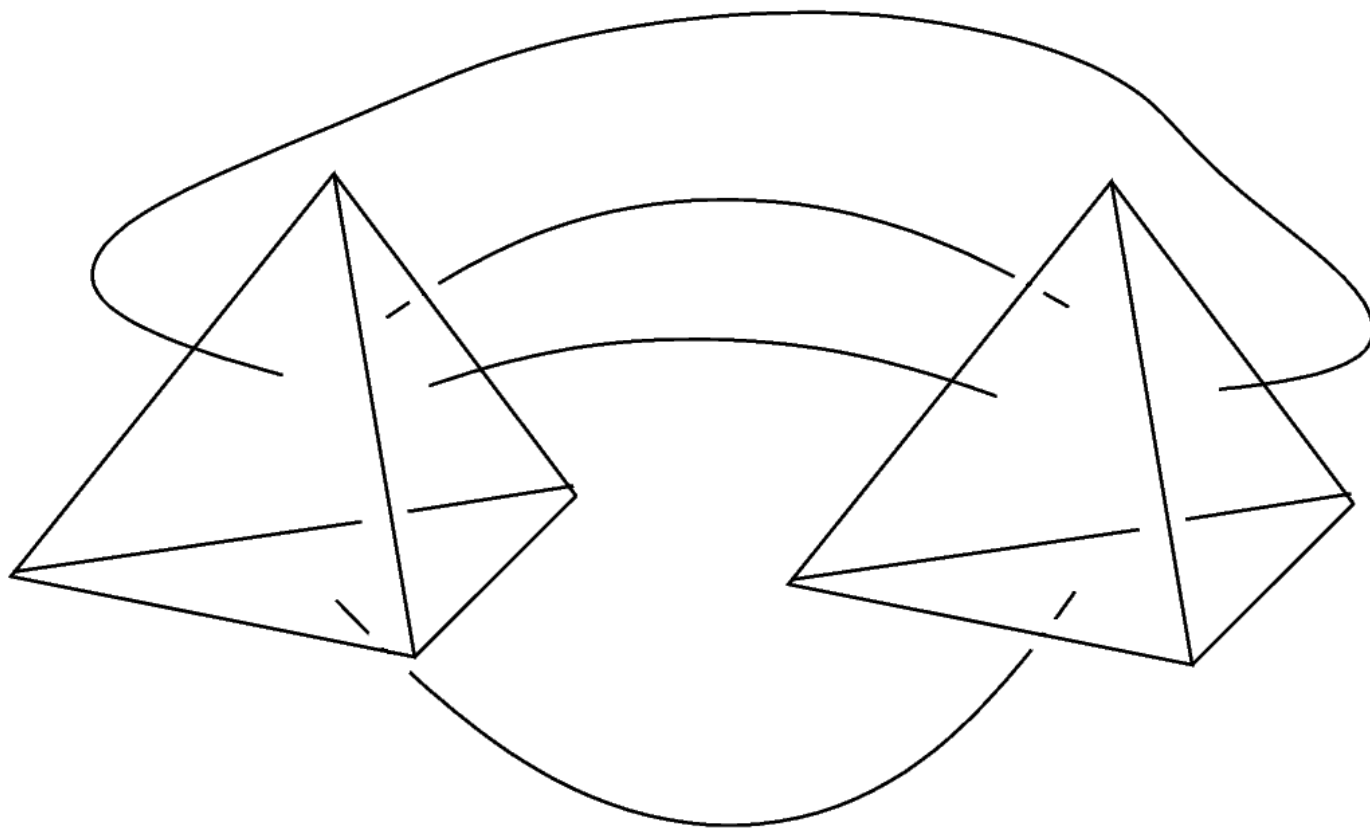
[Joint work with Pawel Wocjan]

Permutational Algorithms for 3-manifold invariants

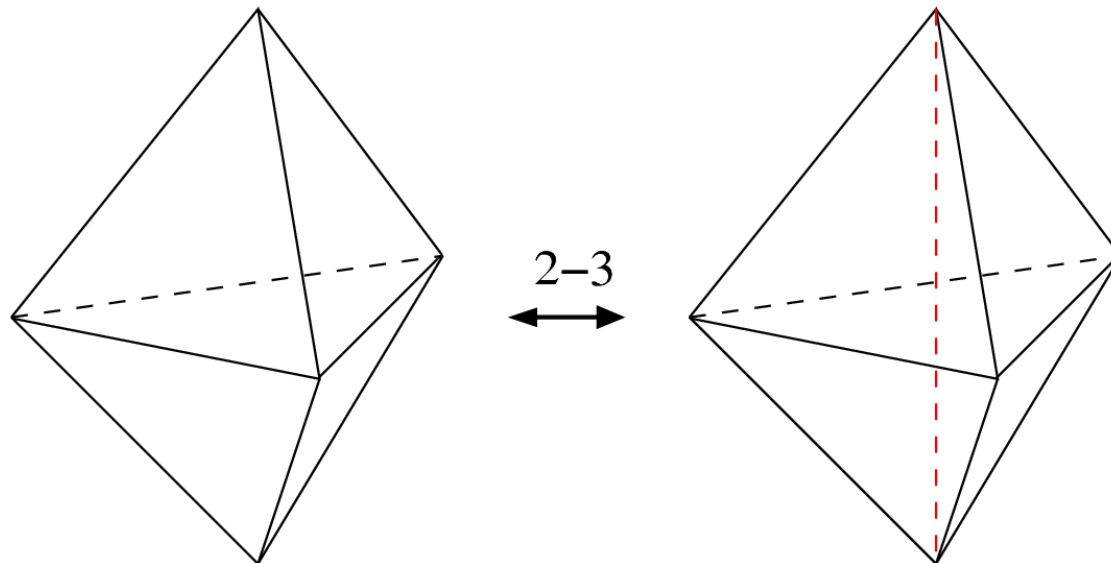
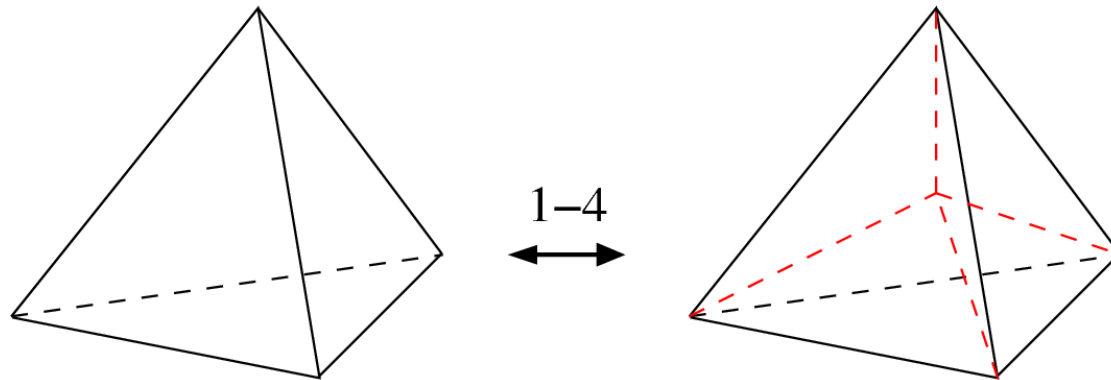
- **3-manifold**: topological space locally like \mathbb{R}^3
- **homeomorphism**: a bijective, continuous map between manifolds whose inverse is continuous
- if a homeomorphism exists between a pair of manifolds we consider them equivalent



- How do we describe a 3-manifold to a computer?
- one way is to use a **triangulation**:
 - a set of tetrahedra
 - a gluing of the faces



- two triangulations yield equivalent 3-manifolds iff they are connected by a finite sequence of Pachner moves



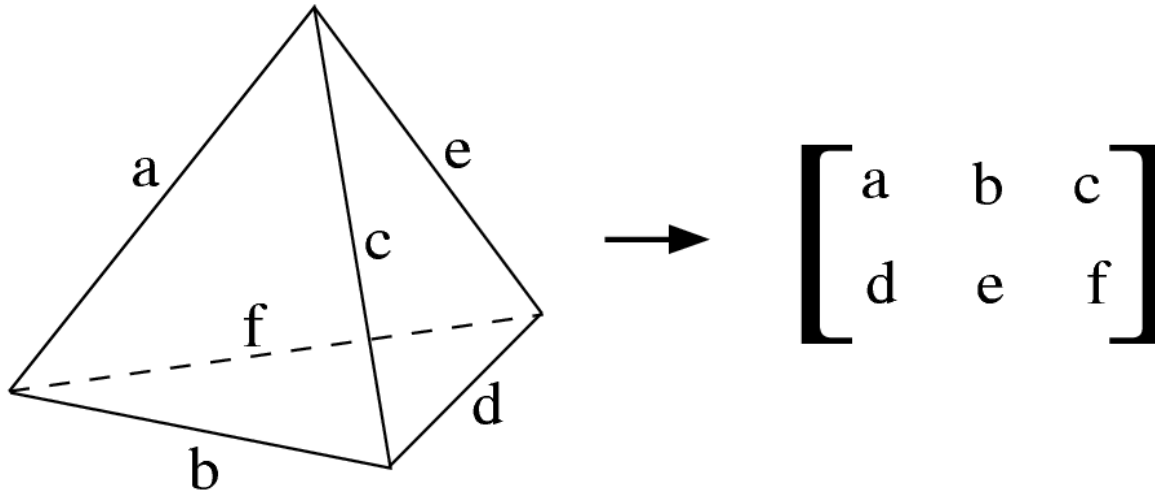
- deciding equivalence of manifolds is not easy!

	equivalence
2-manifolds	in P
3-manifolds	computable
4-manifolds	uncomputable

- partial solution:
manifold invariant – if manifolds A and B are diffeomorphic then $f(A) = f(B)$

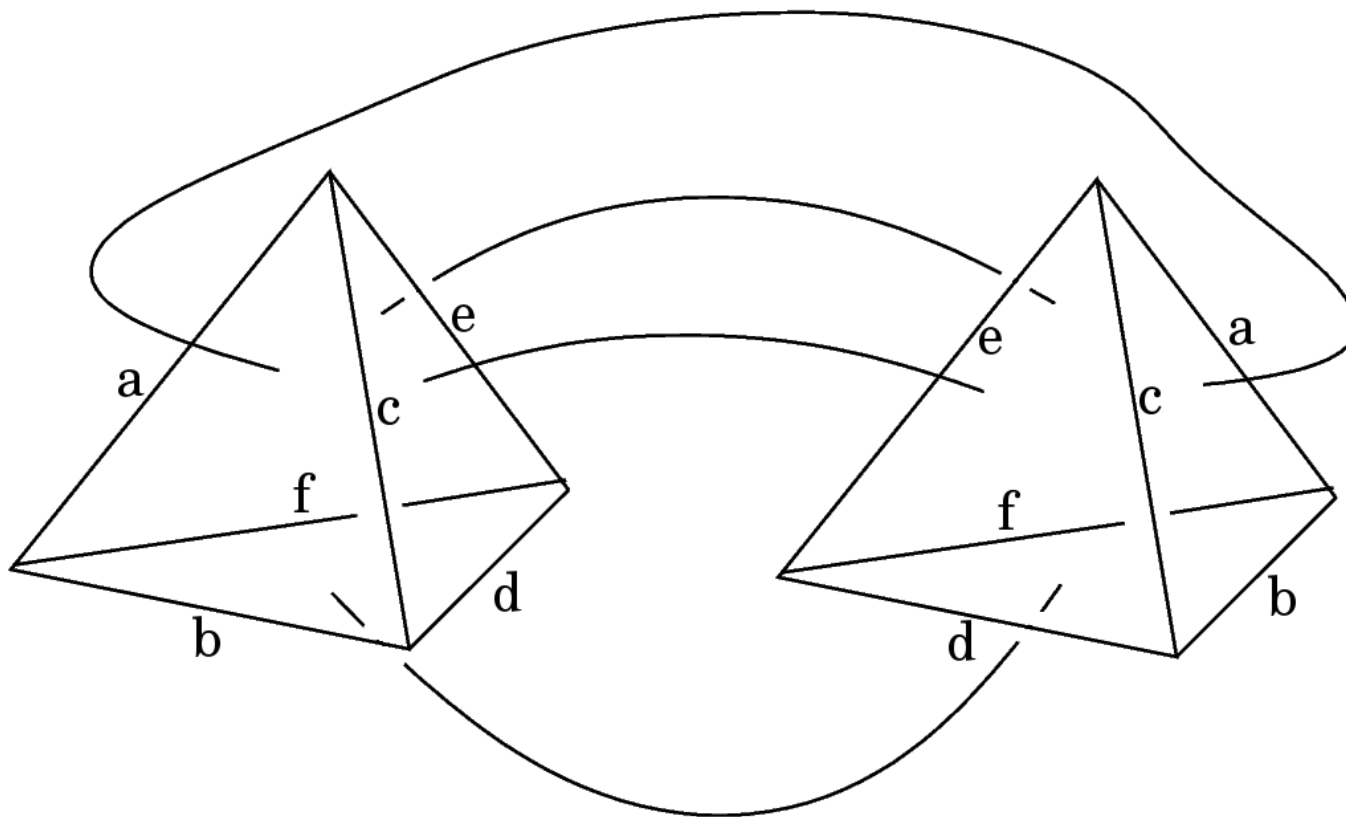
Ponzano-Regge Invariant

- to each tetrahedron, associate one recoupling tensor (one index to each edge)

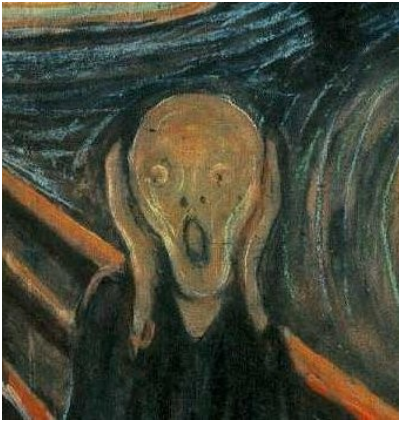


- for each glued face, contract (sum over) the corresponding indices

Example



$$\sum_{abcdef} \begin{bmatrix} a & b & c \\ d & e & f \end{bmatrix} \begin{bmatrix} c & b & a \\ f & e & d \end{bmatrix}$$

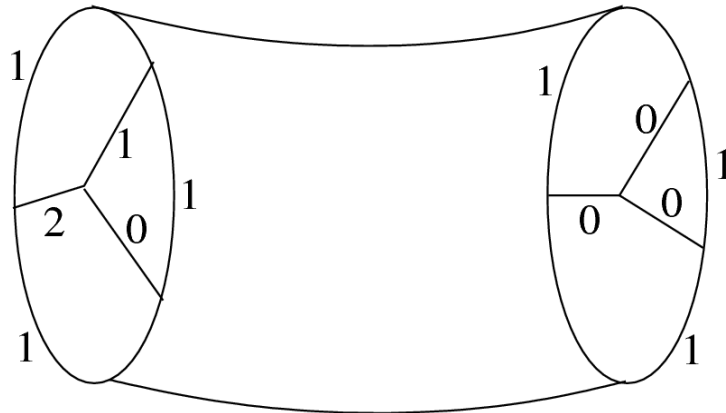


Divergences



- For some triangulations the Ponzano-Regge tensor network is an infinite sum
- Any pair of triangulations of a given manifold such that the sum has finitely many terms yield the same value for the Ponzano-Regge invariant

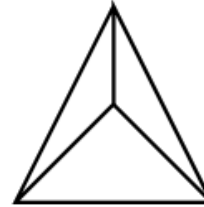
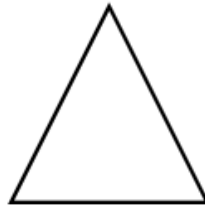
- Boundary is triangulated surface
- we have j labels on edges of triangulation
- these specify a geometry



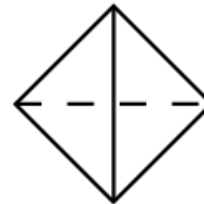
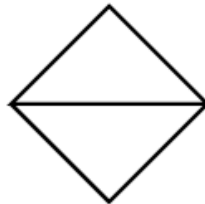
- the value of Ponzano-Regge tensor network is a transition amplitude between geometries
- sum over cobordisms, obtain model of topological quantum gravity

- Gluing of tetrahedra induces a change of triangulation:

subdivide

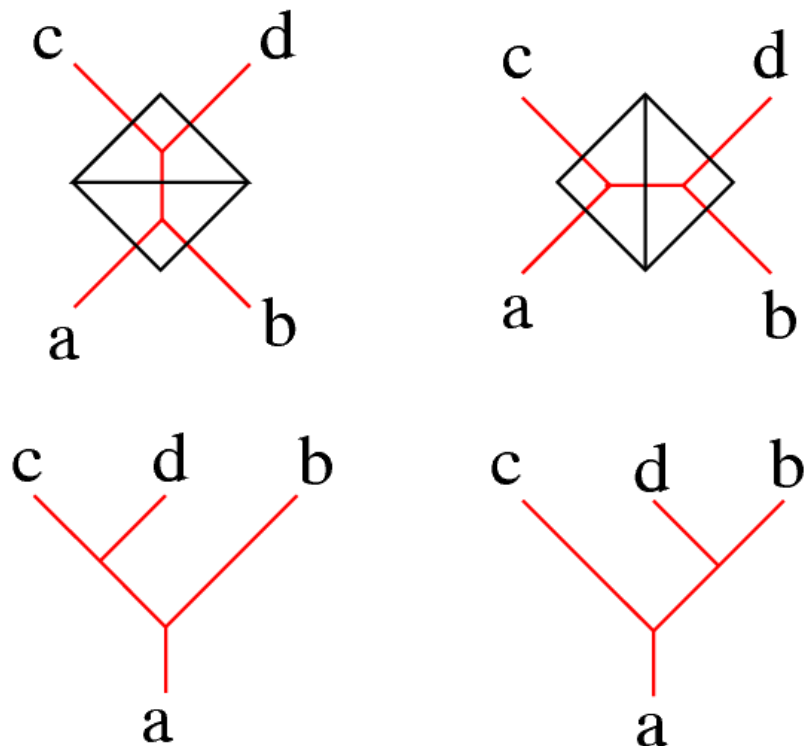


flip



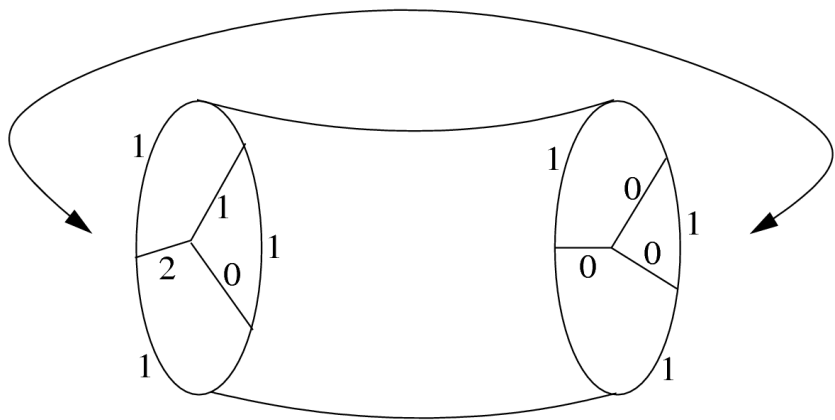
- The Ponzano-Regge amplitude corresponding to the flip move is approximable in PQP

- flip move is **F** move on the dual:

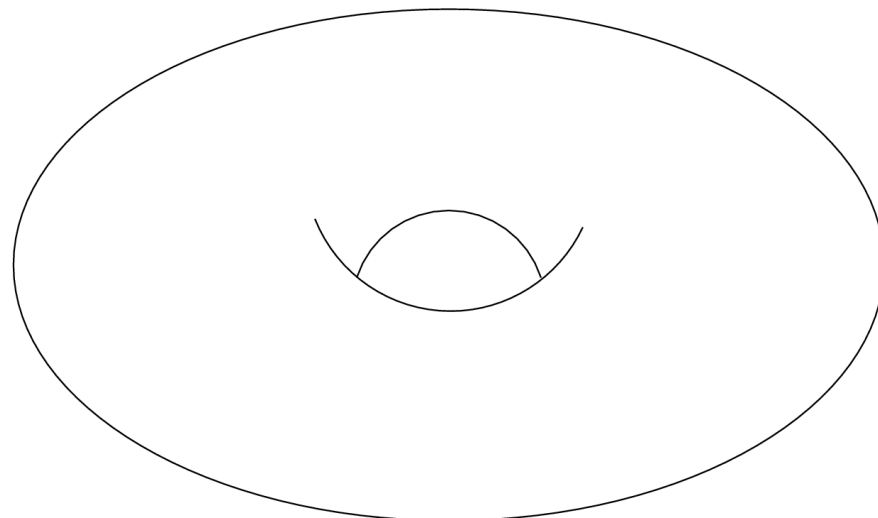


- Hence, in PQP we approximate Ponzano-Regge amplitude for 3-manifolds such that:
 - dual triangulation of boundaries are trees
 - tetrahedra glued two-faces at a time (flip moves)

glue



$$\langle a | U | b \rangle$$



mapping torus



$$\text{Tr}[U]$$

One Clean Qubit

	B_n	S_n
Matrix Elements	BQP-complete	\subset BQP
Characters	DQC1-complete	\subset BPP

Normalized Characters of S_n in BPP Proof:

Theorem 1 (Roichman) For any partitions $\mu = (\mu_1, \dots, \mu_l)$ and $\lambda = (\lambda_1, \dots, \lambda_k)$ of n , the corresponding irreducible character of S_n is given by

$$\chi_\mu^\lambda = \sum_{\Lambda} W_\mu(\Lambda)$$

where the sum is over all standard Young tableaux Λ of shape λ and

$$W_\mu(\Lambda) = \prod_{\substack{1 \leq i \leq k \\ i \notin B(\mu)}} f_\mu(i, \Lambda)$$

where $B(\mu) = \{\mu_1 + \dots + \mu_r \mid 1 \leq r \leq l\}$ and

$$f_\mu(i, \Lambda) = \begin{cases} -1 & \text{box } i+1 \text{ of } \Lambda \text{ is in the southwest of box } i \\ 0 & \text{ } i+1 \text{ is northeast of } i, i+2 \text{ is southwest of } i+1, \text{ and } i+1 \notin B(\mu) \\ 1 & \text{otherwise} \end{cases}$$

Theorem 2 (Greene, Nijenhuis, and Wilf) With polynomial resources, one can sample uniformly from the standard Young Tableaux corresponding to a given shape (n -box Young diagram) using the Hook walk algorithm.



End of Proof.

Is it Universal?

- So far we have seen:
 - $PQP \subset BQP$
 - probably $PQP \not\subseteq P$
- $PQP = BQP$?
- My intuition: No
 - no density
 - one clean qubit version is in BPP

Fault Tolerance

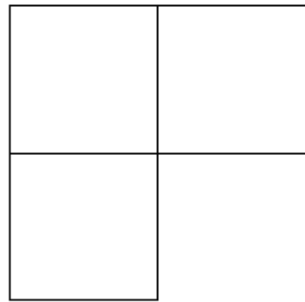
- Model is inherently discrete
 - like topological model
 - unlike circuit model
- Computation occurs in a noiseless subsystem for uniform magnetic fields
 - total angular momentum operators commute with magnetic field operators

Some Open Questions

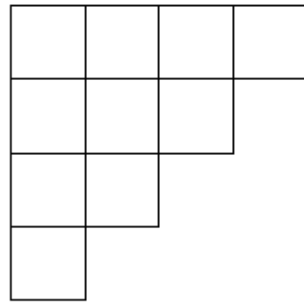
- Implementation of phase measurements?
- q -deformed version of Ponzano-Regge algorithm (Turaev-Viro)
- How computationally powerful are spin-foam models for general topologies?

- For an exponentially large unitary matrix the average magnitude of the matrix elements is exponentially small.
- We approximate to polynomial precision?
- Is this trivial?
 - For random instances: yes.
 - In worst case: probably not.

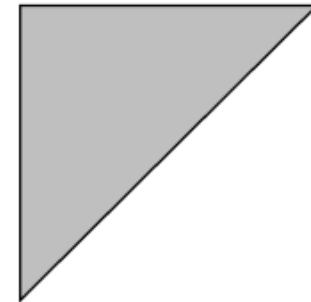
- The normalized character tells us the average diagonal element.
- In certain cases this is large.



3



10



∞

$$\frac{\chi_{\lambda_n}(\pi)}{d_{\lambda_n}} = C_{\pi}(\omega) n^{-|\pi|/2} + O(n^{-|\pi|/2-1})$$

Young's Orthogonal Form

$$\rho_\lambda(\sigma_i)\Lambda = \frac{1}{\tau_i^\Lambda}\Lambda + \sqrt{1 - \frac{1}{(\tau_i^\Lambda)^2}}\Lambda'$$

$$\rho_{\begin{array}{|c|} \hline \square \\ \hline \square \\ \hline \end{array}}(\sigma_2) \begin{array}{|c|c|} \hline 1 & 2 \\ \hline 3 & \\ \hline \end{array} = -\frac{1}{2} \begin{array}{|c|c|} \hline 1 & 2 \\ \hline 3 & \\ \hline \end{array} + \frac{\sqrt{3}}{2} \begin{array}{|c|c|} \hline 1 & 3 \\ \hline 2 & \\ \hline \end{array}$$

1	2	4
3	5	6

