

Quantum Processes and Computation

Assignment 2, Monday 28 Oct 2024

Solutions are shown after each question. Note some solutions are marked *Sketch*. These are intended to be instructions on how to work out the solution yourself, rather than an example of how you should answer this question on an exam.

Exercise 1: We can write the cup/cap for any dimension as a sum over ONB elements:

$$\cup = \sum_{i=1}^d \downarrow_i \downarrow_i \qquad \cap = \sum_{i=1}^d \uparrow_i \uparrow_i$$

(i) Using this definition (and not the matrix form) verify the yanking equations.

$$\text{cup with wire} = \text{wire} \qquad \text{cap with wire} = \text{cup}$$

(ii) Compute the matrices for the cup and cap in 3 dimensions.

Begin Solution:

(i) These can be verified using the properties of ONBs and sums. For the first one:

$$\begin{aligned} \text{cup with wire} &= \sum_i \uparrow_i \uparrow_i \downarrow_j = \sum_{ij} \uparrow_i \downarrow_j \downarrow_j \\ &= \sum_{ij} \delta_{ij} \downarrow_i = \sum_i \downarrow_i = \text{wire} \end{aligned}$$

For the second one:

$$\text{cap with wire} = \sum_i \downarrow_i \downarrow_i = \sum_i \downarrow_i \downarrow_i = \text{cup}$$

(ii) This is a column vector, whose entries correspond to the basis elements $|i, j\rangle$ for $i, j \in \{0, 1, 2\}$. If $i = j$, the entry is a 1, otherwise it is a 0. This results in the following vector:

$$\begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \end{pmatrix}$$

n.b. it is the same as the columns of a 3×3 identity matrix, all stacked on top of each other. Cups in all dimensions have this same pattern. The cap is the transpose of the above, which is row vector:

$$(1 \ 0 \ 0 \ 0 \ 1 \ 0 \ 0 \ 0 \ 1)$$

End Solution

Exercise 2 (5.86): This exercise is about encoding classical functions as linear maps using ONB states and effects, as explained in Section 5.3.4. For a function $F : \{0, 1\}^m \rightarrow \{0, 1\}^n$, we can define an associated linear map f as follows:

$$f = \sum_{(a_1 \dots a_m \mapsto b_1 \dots b_n) \in F} \begin{array}{c} \downarrow b_1 \quad \dots \quad \downarrow b_n \\ \uparrow a_1 \quad \dots \quad \uparrow a_m \end{array}$$

where the notation $(a_1 \dots a_m \mapsto b_1 \dots b_n) \in F$ means we are summing over the *graph* of F , i.e. the set of bitstrings $\{(a_1, \dots, a_m, b_1, \dots, b_n) \mid F(a_1, \dots, a_m) = (b_1, \dots, b_n)\}$.

Using this encoding, define:

$$\begin{aligned} \text{XOR} &= \begin{array}{c} \downarrow 0 \\ \uparrow 0 \ \uparrow 0 \end{array} + \begin{array}{c} \downarrow 1 \\ \uparrow 0 \ \uparrow 1 \end{array} + \begin{array}{c} \downarrow 1 \\ \uparrow 1 \ \uparrow 0 \end{array} + \begin{array}{c} \downarrow 0 \\ \uparrow 1 \ \uparrow 1 \end{array} \\ \text{CNOT} &:= \begin{array}{c} \downarrow 0 \ \downarrow 0 \\ \uparrow 0 \ \uparrow 0 \end{array} + \begin{array}{c} \downarrow 0 \ \downarrow 1 \\ \uparrow 0 \ \uparrow 1 \end{array} + \begin{array}{c} \downarrow 1 \ \downarrow 1 \\ \uparrow 1 \ \uparrow 0 \end{array} + \begin{array}{c} \downarrow 1 \ \downarrow 0 \\ \uparrow 1 \ \uparrow 1 \end{array} \\ \text{COPY} &:= \begin{array}{c} \downarrow 0 \ \downarrow 0 \\ \uparrow 0 \end{array} + \begin{array}{c} \downarrow 1 \ \downarrow 1 \\ \uparrow 1 \end{array} \end{aligned}$$

Show that

$$\text{CNOT} = \begin{array}{c} \text{XOR} \\ \text{COPY} \end{array}$$

(Hint: try comparing the LHS to the RHS on all basis states, rather than writing out a big sum.)

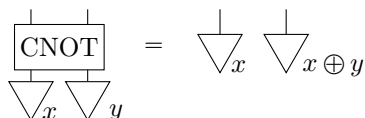
Next, find ψ and ϕ such that the following equation holds:

$$\begin{array}{c} \text{XOR} \\ \text{COPY} \end{array} = \begin{array}{c} \downarrow \phi \\ \uparrow \psi \end{array}$$

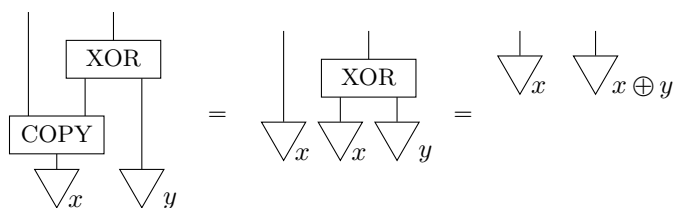
Although it might not look like much now, this equation will turn out to lie at the heart of the notion of *complementarity* which is an important part of the ZX-calculus.

Begin Solution:

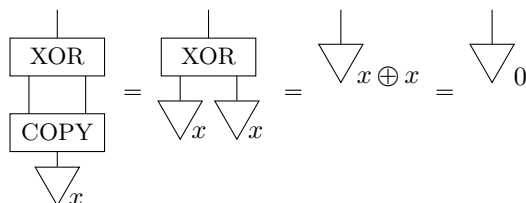
The first part can be done by plugging in basis states, and nothing that the LHS gives:



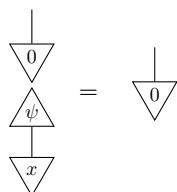
and the RHS gives:



If I evaluate the second diagram at a basis state, I get:

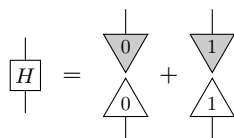


This tells me that $|\phi\rangle$ should be $|0\rangle$. For $\langle\psi|$, I need the effect that “deletes” any basis state: $\langle\psi| = \langle 0| + \langle 1| = \sum_i \langle i|$. Then $\langle\psi|0\rangle = \langle\psi|1\rangle = 1$, so:



End Solution

Exercise 3: Let the Hadamard gate, which sends the Z-basis to the X-basis be defined as follows:



where

$$\begin{array}{c} \downarrow \\ 0 \end{array} := \frac{1}{\sqrt{2}} \left(\begin{array}{c} \downarrow \\ 0 \end{array} + \begin{array}{c} \downarrow \\ 1 \end{array} \right) \quad \begin{array}{c} \downarrow \\ 1 \end{array} := \frac{1}{\sqrt{2}} \left(\begin{array}{c} \downarrow \\ 0 \end{array} - \begin{array}{c} \downarrow \\ 1 \end{array} \right)$$

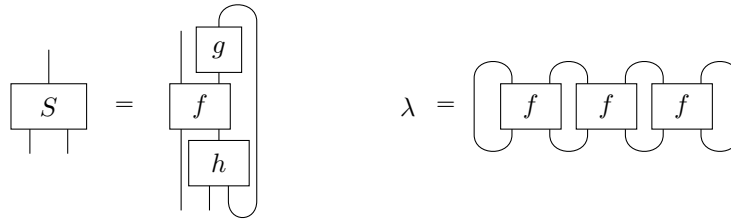
Compute the matrix of H . Show that $H = H^\dagger = H^T$. Using this fact (or otherwise) show that H also sends the X-basis back to the Z-basis.

Begin Solution:

Sketch: The matrix can be computed by plugging in each of the 4 bras and kets of the computational basis. We can see it sends X-basis elements to Z-basis elements by applying the adjoint to both sides of the definition of H and using $H = H^\dagger$.

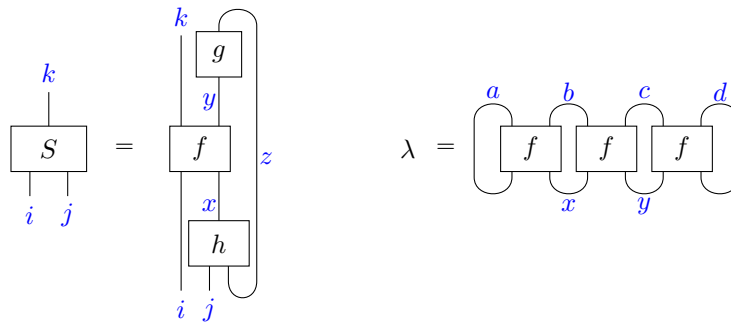
End Solution

Exercise 4: Write the following diagrams as tensor contractions, i.e. as sums over products of matrix elements f_{ij}^{kl} , etc.



Begin Solution:

Labelling the wires with some index names:



...we get:

$$S_{ij}^k = \sum_{xyz} f_{ix}^{ky} g_y^z h_{jz}^x$$

$$\lambda = \sum_{abcdxy} f_{ax}^{ab} f_{xy}^{bc} f_{yd}^{cd}$$

Note λ is a scalar, so all indices on the RHS are summed over.

End Solution